



# B.E. International Program

## Faculty of Economics, Thammasat University



TU152 Fundamental Mathematics

Quiz #1 (Semester 2/2012)

Name.....Student No.....

1. (10 marks) Prove that if  $q \vee \sim(u \vee p)$ ,  $\sim(q \wedge r)$ , and  $\sim u \rightarrow (u \vee r)$ , then  $u \leftrightarrow \sim r$ .

Proof: First, we need to show  $u \rightarrow \sim r$ . Assume  $u$  and show  $\sim r$ .

- 1)  $u$  Assumption
- 2)  $u \rightarrow (u \vee p)$  Law of Addition
- 3)  $u \vee p$  1, 2, Rule of Modus Ponens
- 4)  $q \vee \sim(u \vee p)$  Assumption
- 5)  $q$  3, 4, Disjunctive Syllogism, Law of Double Negation, Rule of Substitution
- 6)  $\sim(q \wedge r)$  Assumption
- 7)  $\sim q \vee \sim r$  6, De Morgan's Law, Rule of Substitution
- 8)  $\sim r$  5, 7, Disjunctive Syllogism

Hence,  $u \rightarrow \sim r$ .

Next, we prove  $\sim r \rightarrow u$  by assuming  $\sim r$  and show  $u$ . In this case we will show by using the method of indirect proof, that is we assume  $u$  to be false and we try to show the contradiction.

Note: Since we assume  $u$  to be false then  $\sim u$  is assumed to be true.

- 1)  $\sim u$  Assumption
- 2)  $\sim r$  Assumption
- 3)  $\sim u \wedge \sim r$  1, 2, Definition of Conjunction
- 4)  $\sim(u \vee r)$  3, De Morgan's Law, Rule of Substitution
- 5)  $\sim u \rightarrow (u \vee r)$  Assumption
- 6)  $u$  4, 5, Rule of Modus Tollens, Law of Double Negation, Rule of Substitution
- 7)  $u \wedge \sim u$  1, 6, Definition of Conjunction

Since  $u \wedge \sim u$  is a contradiction, hence  $\sim u$  is false and that is  $u$  is true. Therefore,  $\sim r \rightarrow u$ .

$\therefore$  We have proved that  $u \leftrightarrow \sim r$ .

2. (10 marks) Show that for every natural number  $n$ ,  $n < 2^n$ .

Proof: Let  $S_n$  be the statement “ $n < 2^n$  where  $n$  is any natural number”.

1)  $S_1$  is “ $1 < 2^1$ ” which is true.

2) Assume  $S_k$  is true and we need to show that  $S_{k+1}$  is true.

That is we assume

$$k < 2^k \text{ for every natural number } k.$$

We need to show that  $(k+1) < 2^{(k+1)}$ .

Since  $k < 2^k$

then  $2 \cdot k < 2 \cdot 2^k$

$$k + k < 2^{k+1}$$

$$k + 1 \leq k + k < 2^{k+1} \quad (\text{Because } 1 \leq k)$$

That is,  $k + 1 < 2^{k+1}$

Hence,  $S_{k+1}$  is true.

Therefore, by Mathematical Induction we conclude that for every natural number  $n$ ,  $n < 2^n$ .

## Laws and Rules for Proving:

**Theorem:** For all statements  $p$ ,  $q$ , and  $r$ , the following statements are tautologies.

1. Law of Identity:  $p \leftrightarrow p$
2. Law of Double Negation:  $p \leftrightarrow \sim(\sim p)$
3. Law of Excluded Middle:  $p \vee \sim p$
4. Law of Contradiction:  $\sim(p \wedge \sim p)$
5. Idempotent Laws:  $(p \wedge p) \leftrightarrow p$  ,  $(p \vee p) \leftrightarrow p$
6. Law of Addition:  $p \rightarrow (p \vee q)$
7. Law of Equivalence:  $[p \leftrightarrow q] \leftrightarrow [(p \rightarrow q) \wedge (q \rightarrow p)]$
8. Law of Contraposition:  $(p \rightarrow q) \leftrightarrow (\sim q \rightarrow \sim p)$
9. Law of (Hypothetical) Syllogism:  $[(p \rightarrow q) \wedge (q \rightarrow r)] \rightarrow [p \rightarrow r]$
10. Commutative Laws:  $(p \wedge q) \leftrightarrow (q \wedge p)$   
 $(p \vee q) \leftrightarrow (q \vee p)$
11. Associative Laws:  $[p \wedge (q \wedge r)] \leftrightarrow [(p \wedge q) \wedge r] \leftrightarrow [p \wedge q \wedge r]$   
 $[p \vee (q \vee r)] \leftrightarrow [(p \vee q) \vee r] \leftrightarrow [p \vee q \vee r]$
12. Distributive Laws:  $[p \wedge (q \vee r)] \leftrightarrow [(p \wedge q) \vee (p \wedge r)]$   
 $[p \vee (q \wedge r)] \leftrightarrow [(p \vee q) \wedge (p \vee r)]$
13. Absorption Laws:  $[p \wedge (p \vee q)] \leftrightarrow p$   
 $[p \vee (p \wedge q)] \leftrightarrow p$
14. De Morgan's Laws:  $\sim(p \wedge q) \leftrightarrow (\sim p \vee \sim q)$   
 $\sim(p \vee q) \leftrightarrow (\sim p \wedge \sim q)$
15. Law of Implication:  $(p \rightarrow q) \leftrightarrow (\sim p \vee q)$   
 $(p \vee q) \leftrightarrow (\sim p \rightarrow q)$
16. Law of Negation for Implication:  $\sim(p \rightarrow q) \leftrightarrow (p \wedge \sim q)$

## Rules for Proving:

### 1. Rule of Modus Ponens or Rule of Implication

If  $p \rightarrow q$  and  $p$ , then  $q$

### 2. Rule of Substitution

2.1 Substitute the statement that is tautology (or contradiction) with the other statement.

2.2 Substitute the statement with its equivalent statement.

### 3. Definition of Conjunction

3.1 If  $p \wedge q$ , then  $p$  and  $q$

3.2 If  $p$  and  $q$ , then  $p \wedge q$

### 4. Rule of Modus Tollens

If  $p \rightarrow q$  and  $\sim q$ , then  $\sim p$

### 5. Disjunctive Syllogism

5.1 If  $p \vee q$  and  $\sim p$ , then  $q$

5.2 If  $p \vee q$  and  $\sim q$ , then  $p$

### 6. Constructive Dilemma

If  $(p \rightarrow q) \wedge (r \rightarrow s)$  and  $p \vee r$ , then  $q \vee s$